TraceTracker:

Hardware/Software Co-Evaluation for Large-Scale I/O Workload Reconstruction

> Miryeong Kwon, Jie Zhang, Gyuyoung Park, Wonil Choi, David Donofrio, John Shalf, Mahmut Kandemir, and Myoungsoo Jung



Summary

- <u>Motivation</u>: Block traces collected on old systems (a decade ago HDDs) can make system analysis significantly different from the modern systems
- <u>Challenges</u>: Trace reconstruction is high-cost work and inaccurate
 - Application execution requires prohibitive resources; thousands of users, large-scale computing systems (impractical)
 - Overly-simplified methods are excessively imprecise due to lack of runtime contexts
- Goal: Accurately reconstruct workloads without application execution
- **Solutions**: Hardware/Software co-evaluation method
 - 1. Software method: Infer runtime contexts from existing block traces
 - 2. Hardware method: Remaster storage traces; aware of target system
- Evaluation:
 - TraceTracker infer runtime contexts 99%, 96% for number of occurrences and total periods, respectively
 - Apply TraceTracker to 577 open-license block traces



TraceTracker

1. Background of Block Trace

2. Trace Reconstruction Methods

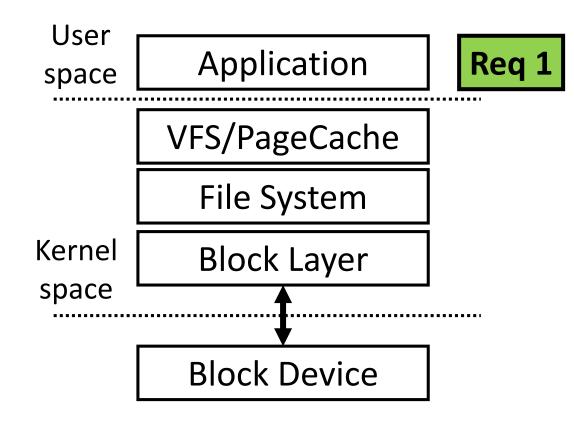
3. Insights on Trace Reconstruction

4. TraceTracker Method

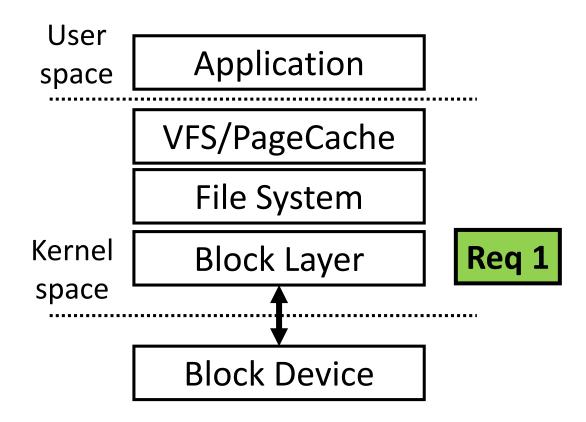
5. Evaluation



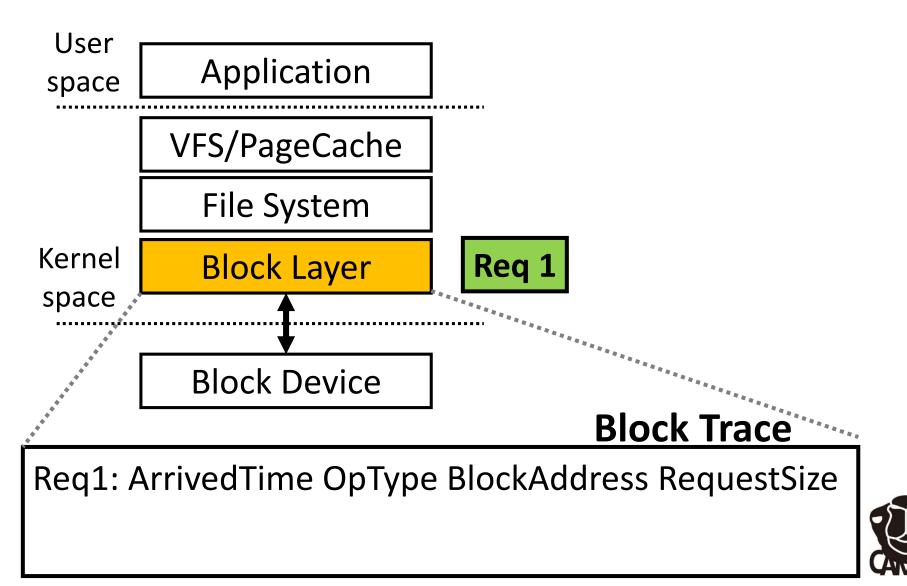


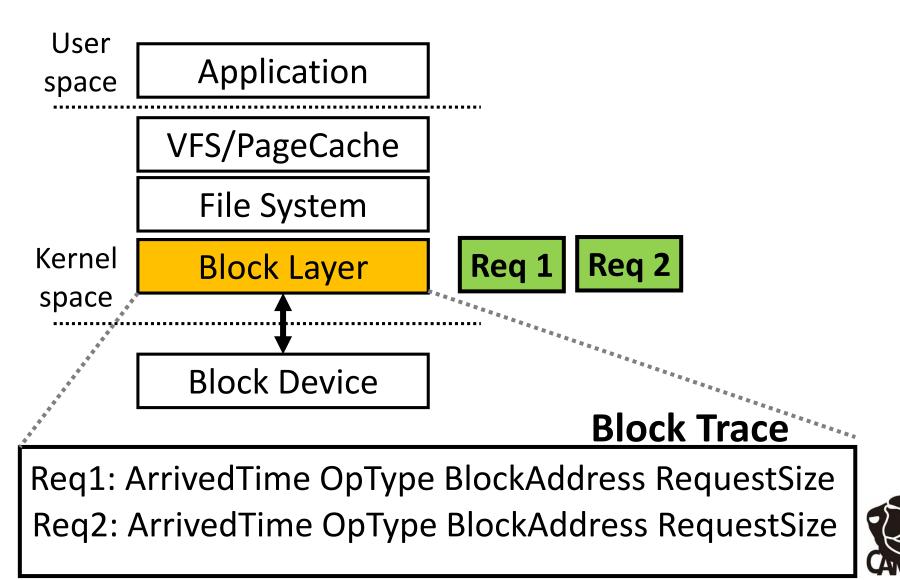


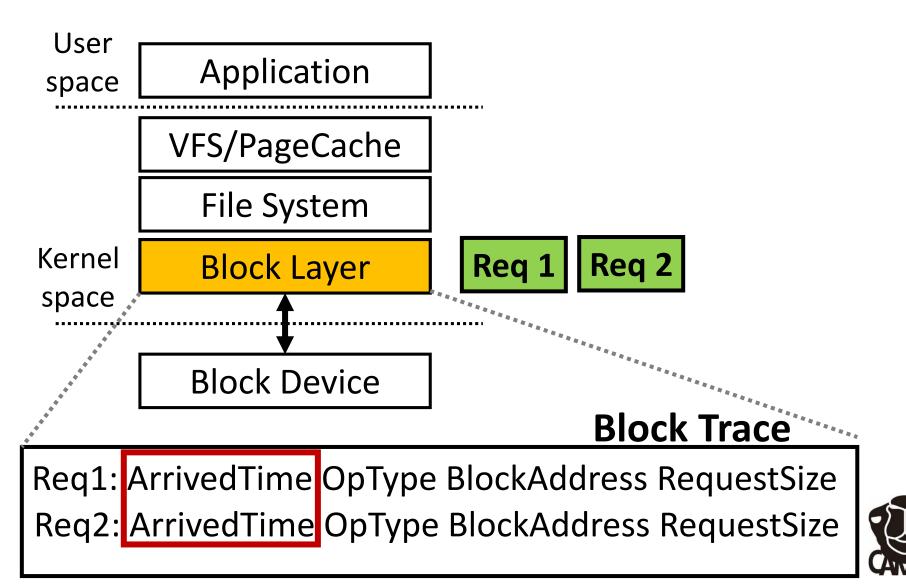




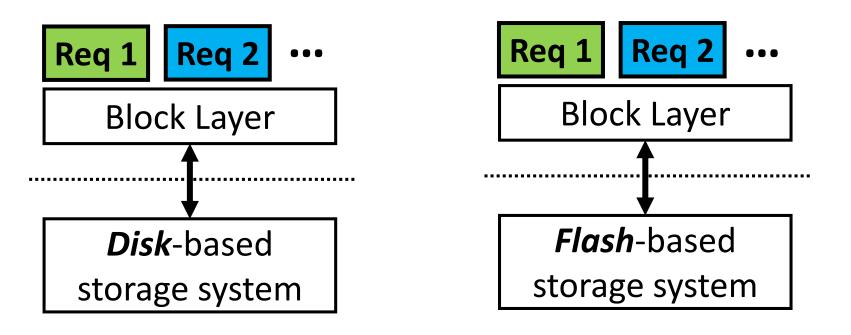




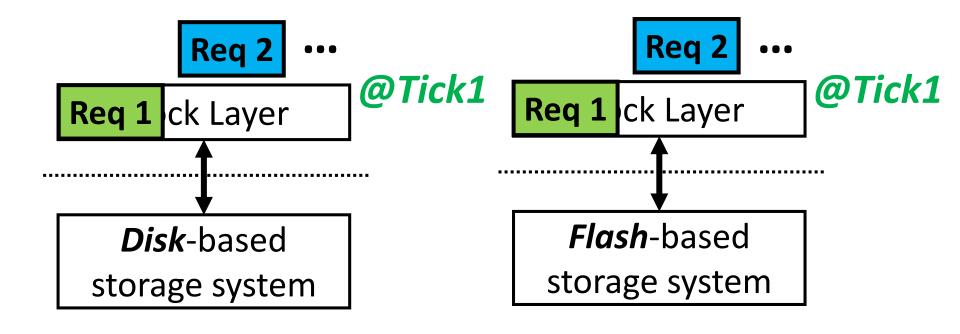




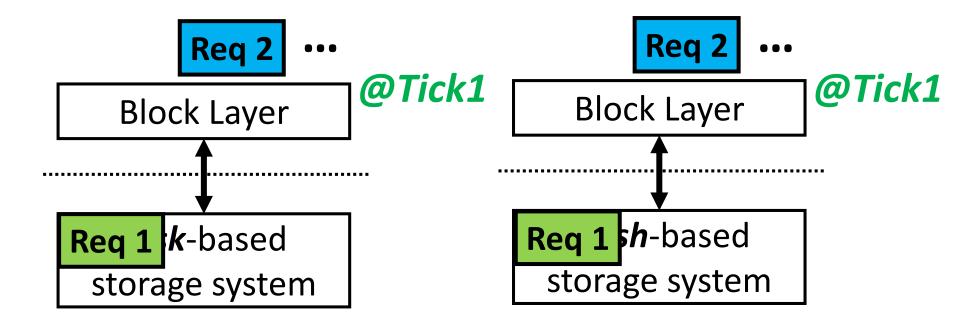




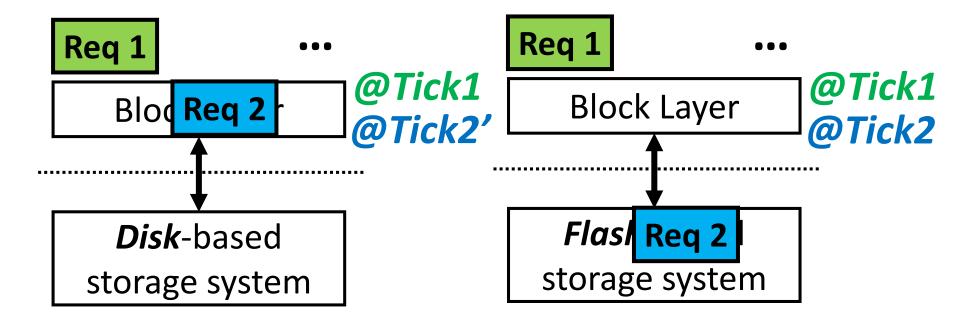






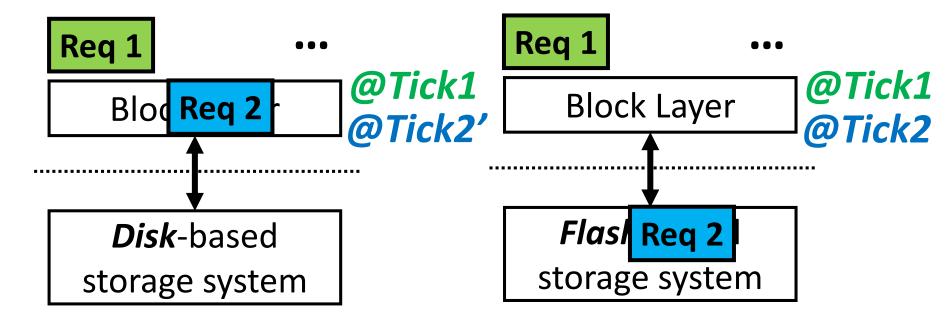








Existing traces: collected on the *disk-based* storage systems and published around **2007**



Timing information of block traces should be revised correspond to target storage system



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• APPLICATION

Large-scale computing system

Target Storage System



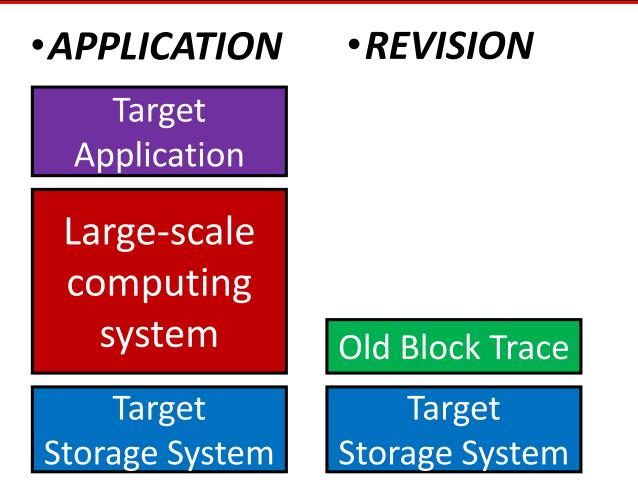
• APPLICATION

Target Application

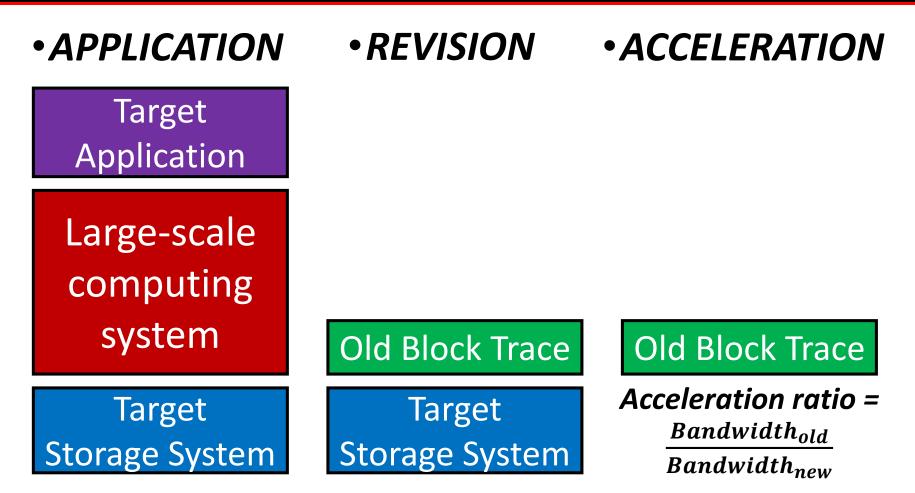
Large-scale computing system

Target Storage System

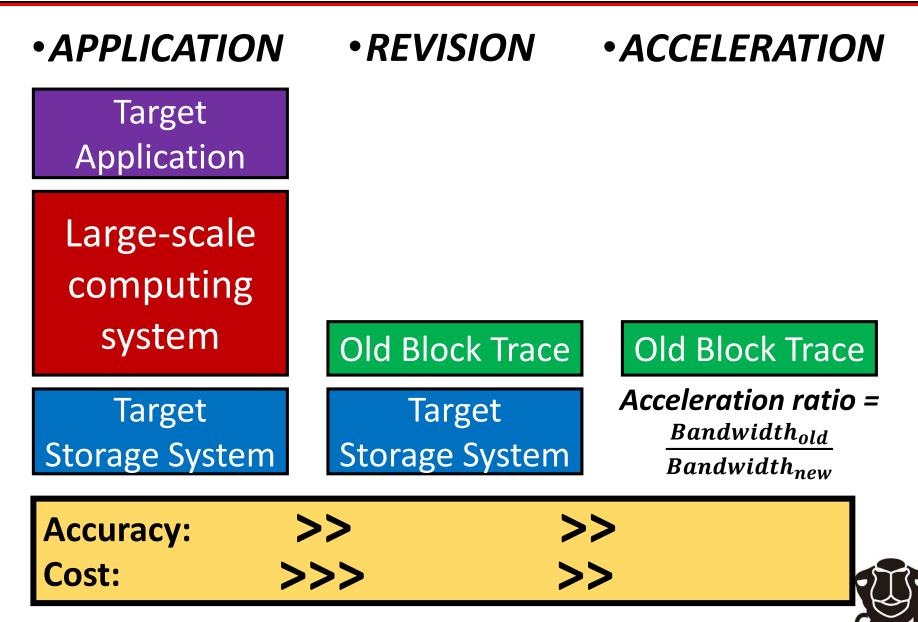


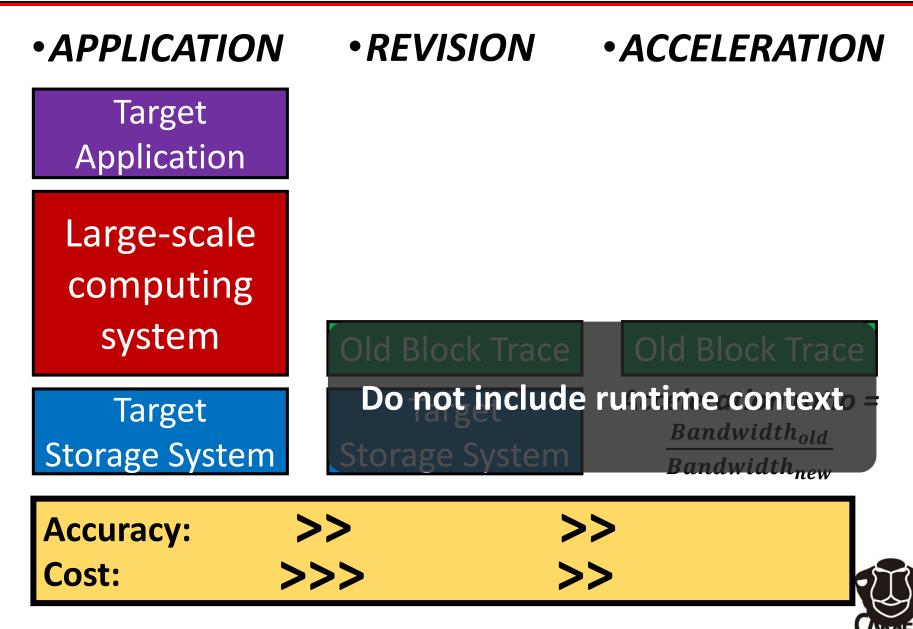


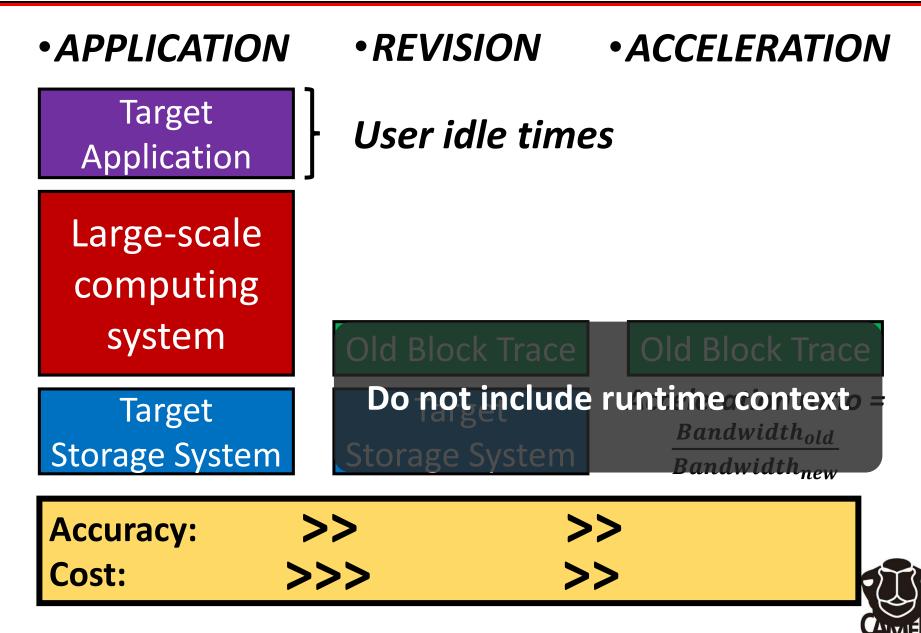


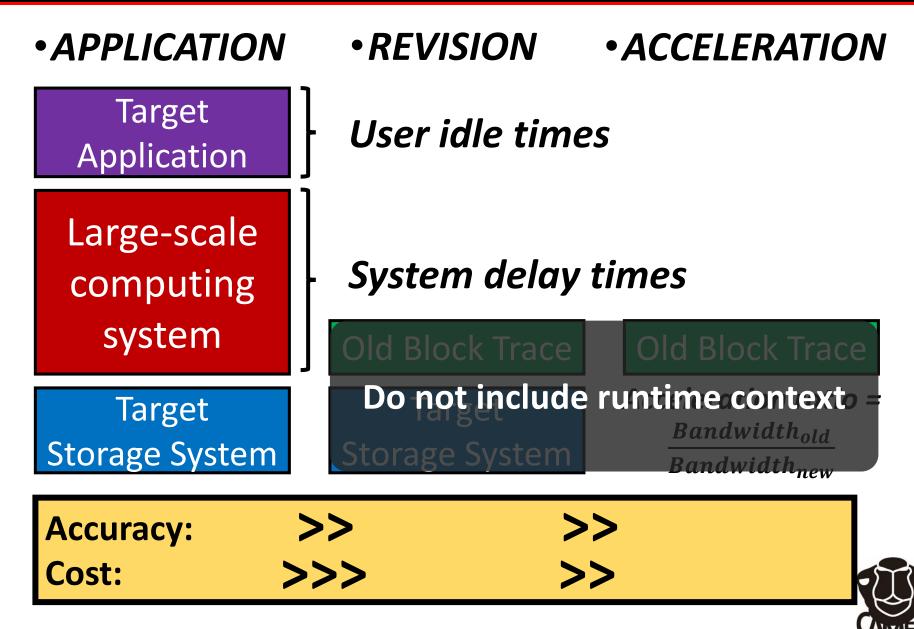


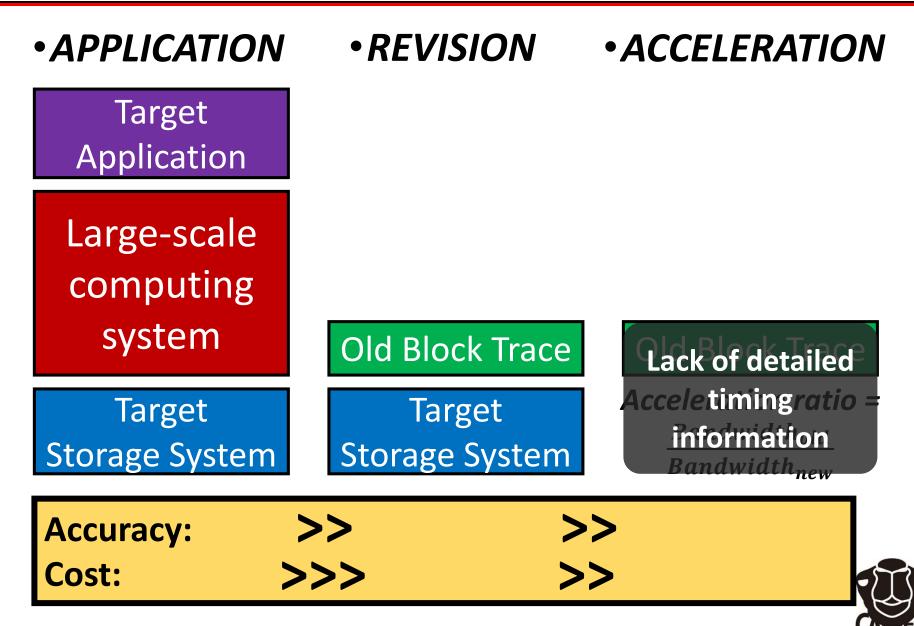


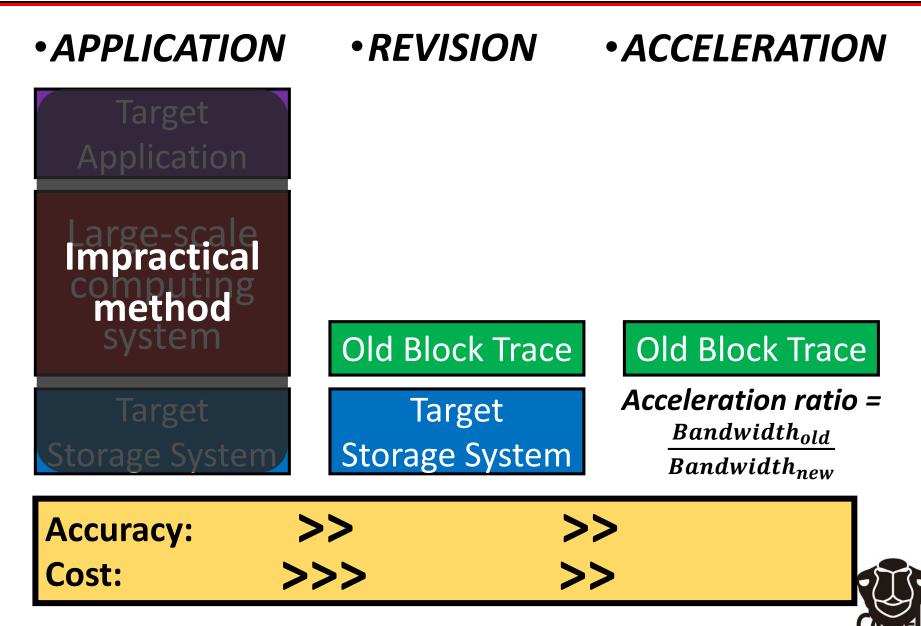


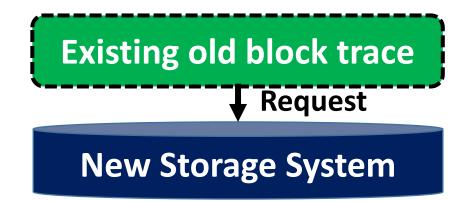




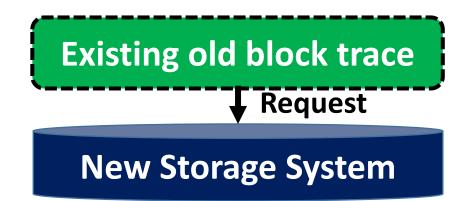






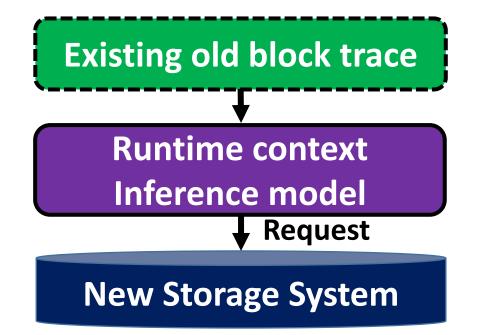






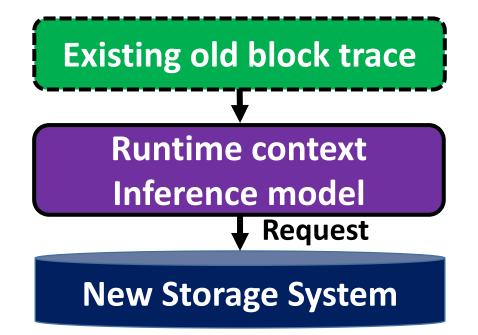
Goal: Accurately reconstruct workloads without application execution (at low cost)





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However, how runtime context can be inferred from block trace?



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Goal: Infer runtime contexts from timing information (Inter-arrival times) of old block traces



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Insight1: Inter-arrival times decomposition



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Insight1: Inter-arrival times decomposition

Insight2: Three inter-arrival time types



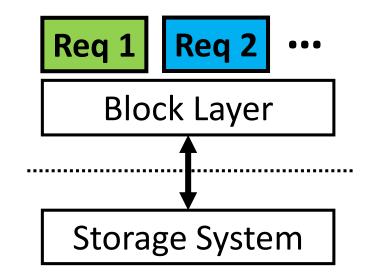
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Insight1: Inter-arrival times decomposition

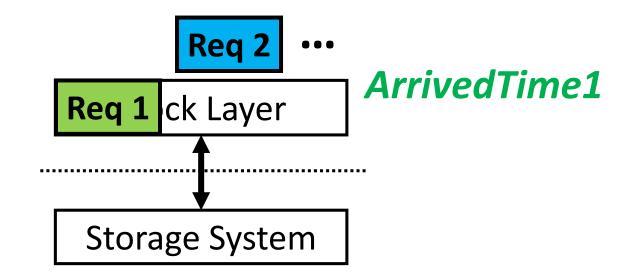
Insight2: Three inter-arrival time types

Insight3: Inter-arrival time components' characteristics

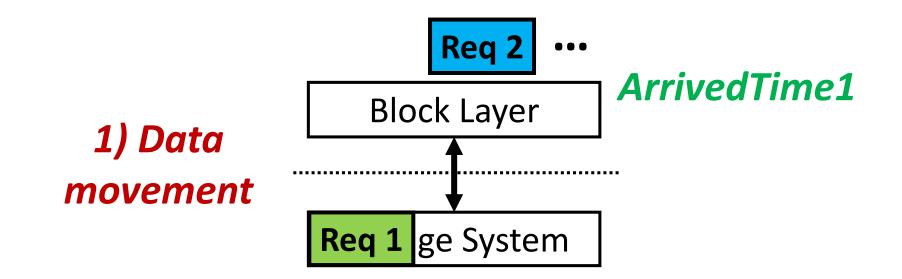




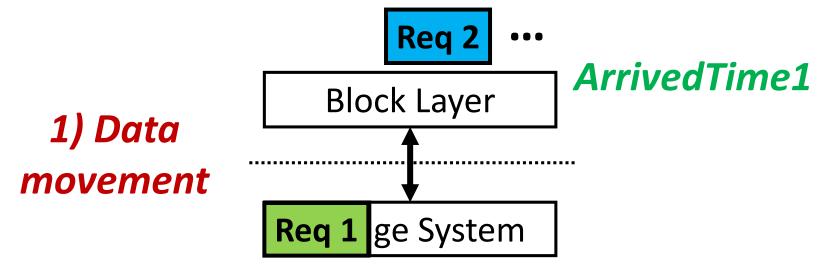




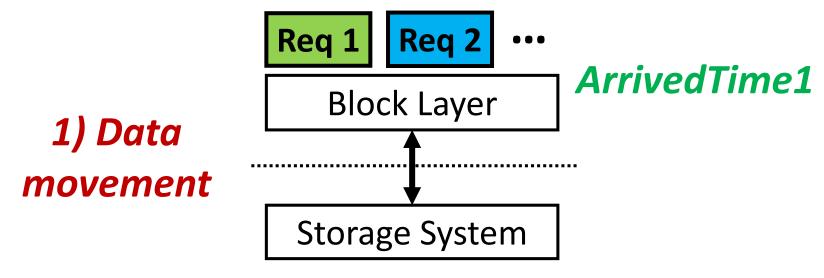






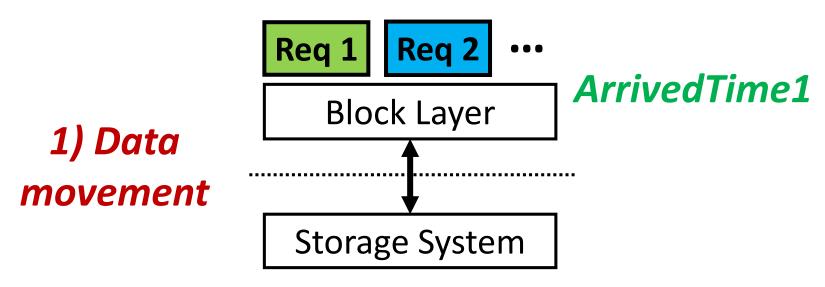






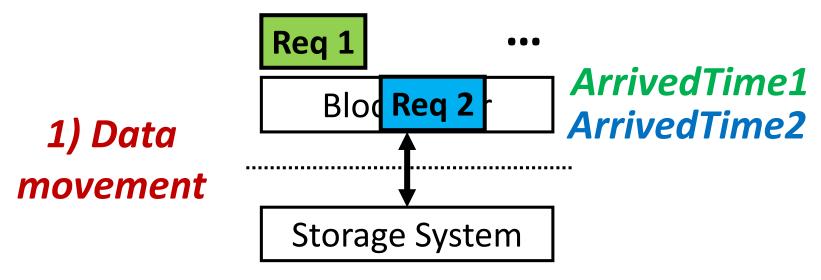


3) User idle times

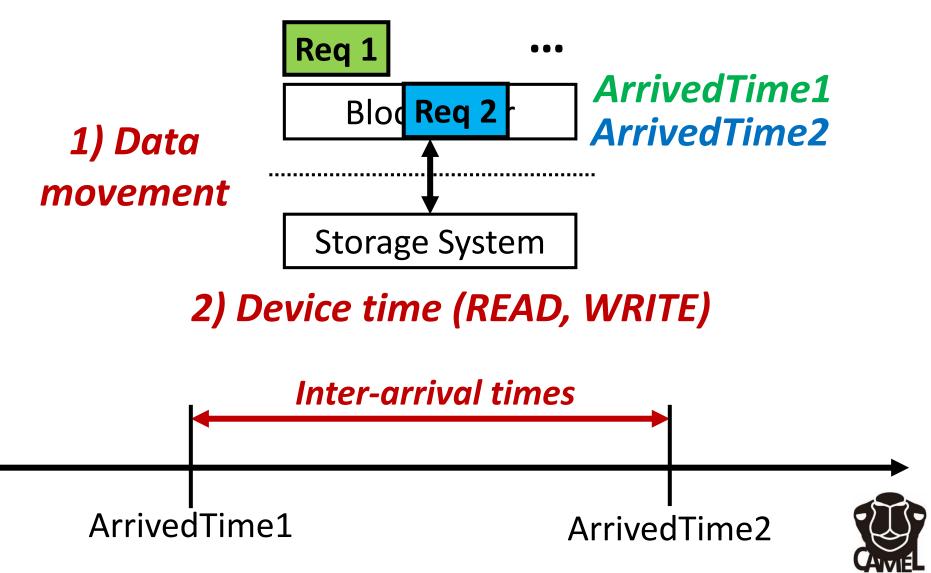


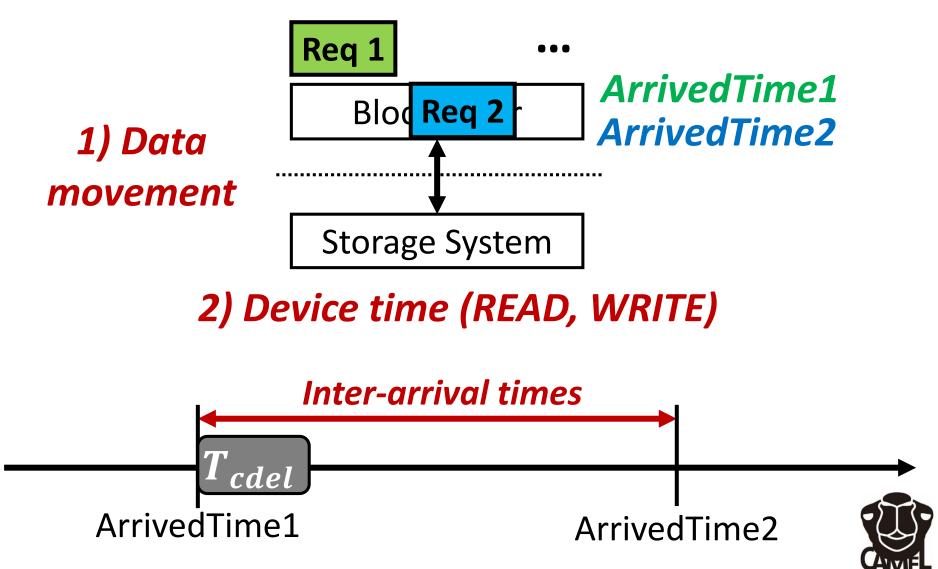


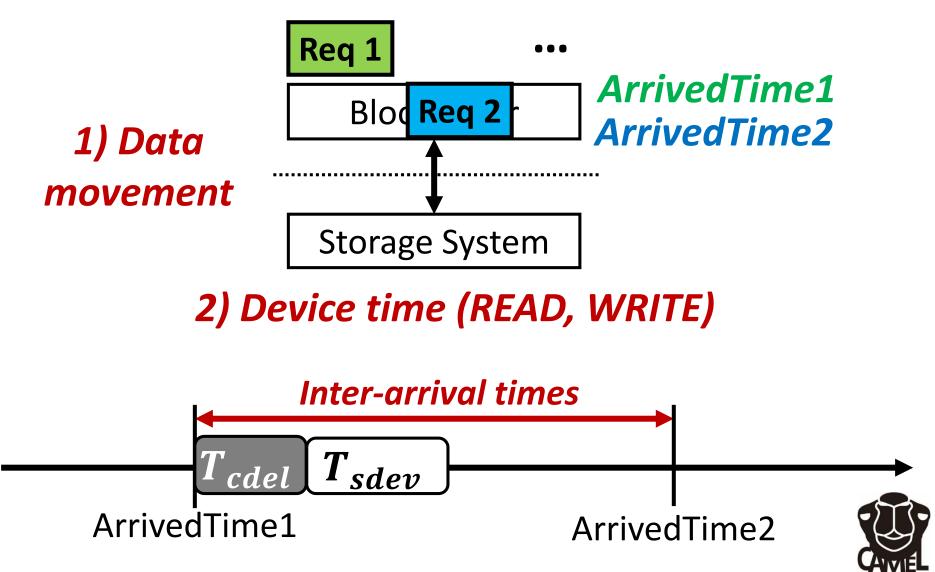
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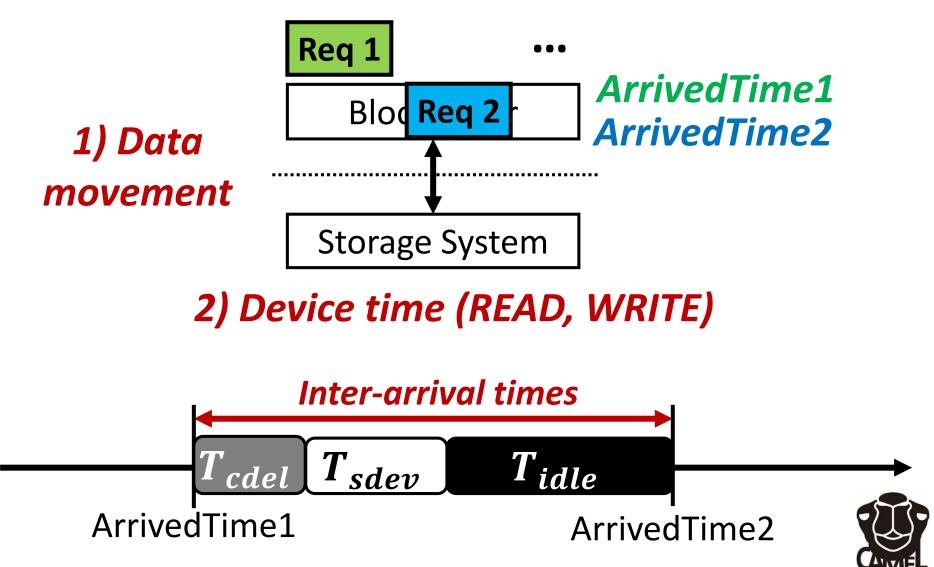














Sync I/O, Non-busy status T_{cdel} T_{sdev} T_{idle}



Sync I/O, Non-busy status



Sync I/O, Busy status





Sync I/O,
Non-busy status T_{cdel} T_{sdev} T_{idle} Sync I/O,
Busy status T_{cdel} T_{sdev}



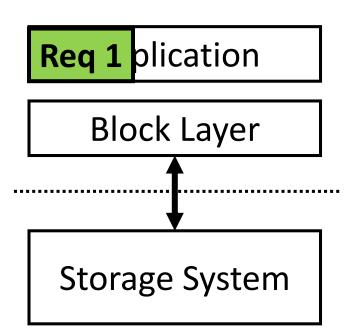




	T _{cdel}	T _{sdev}	T _{idle}
Time length	short		long
Predictable	0	0	X (rare)

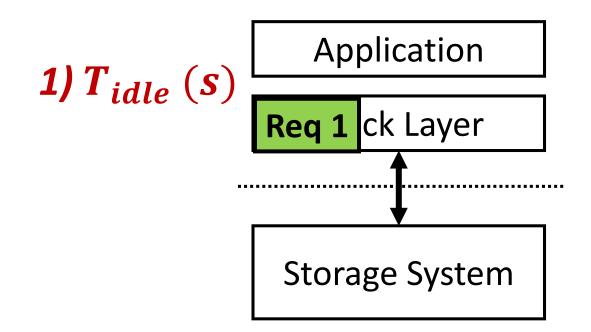


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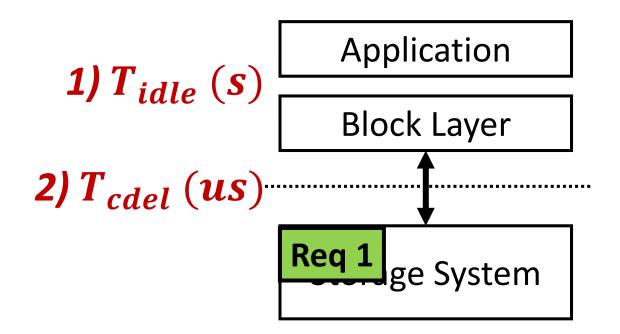


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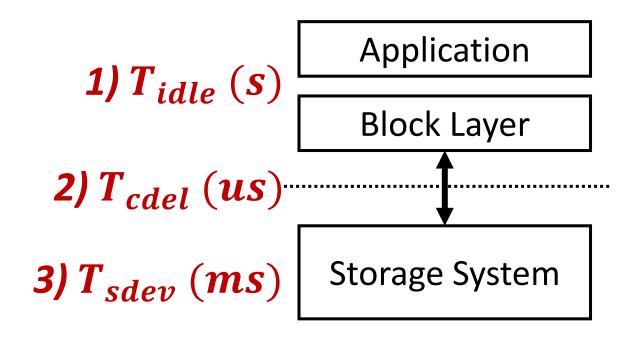


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Pre-software evaluation



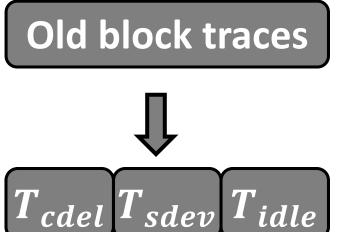






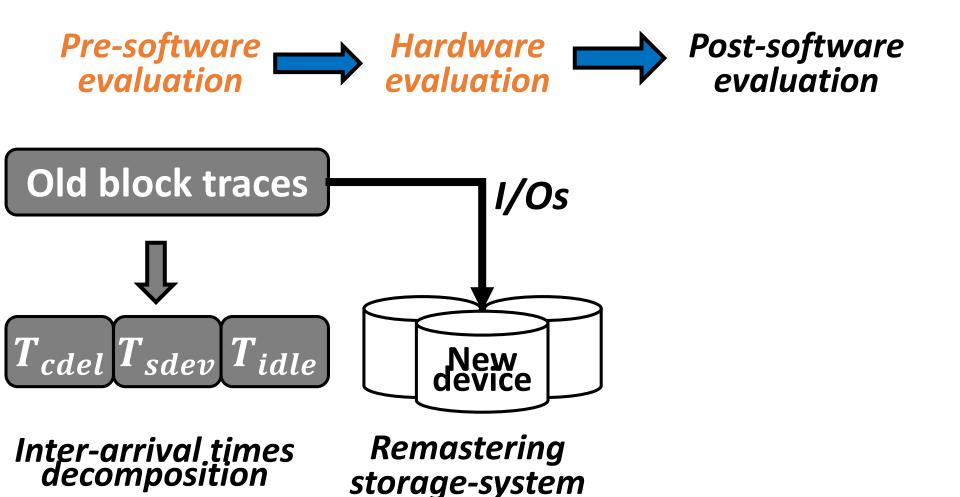




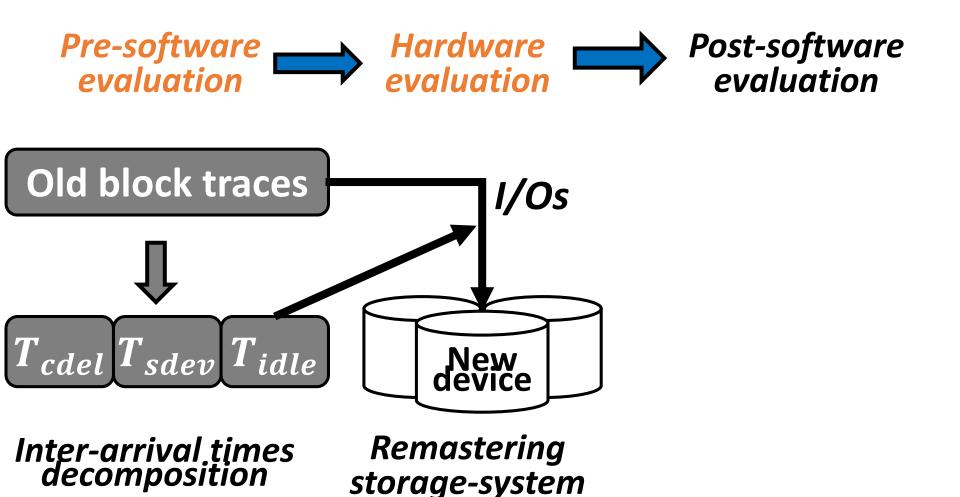


Inter-arrival times decomposition

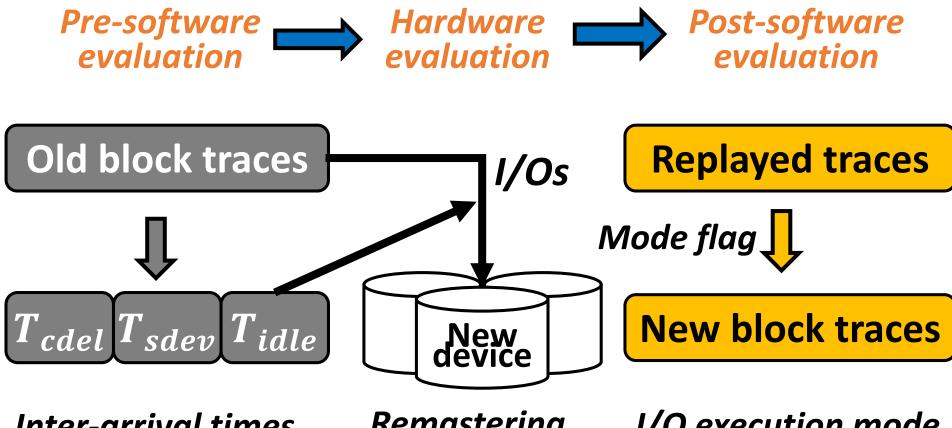








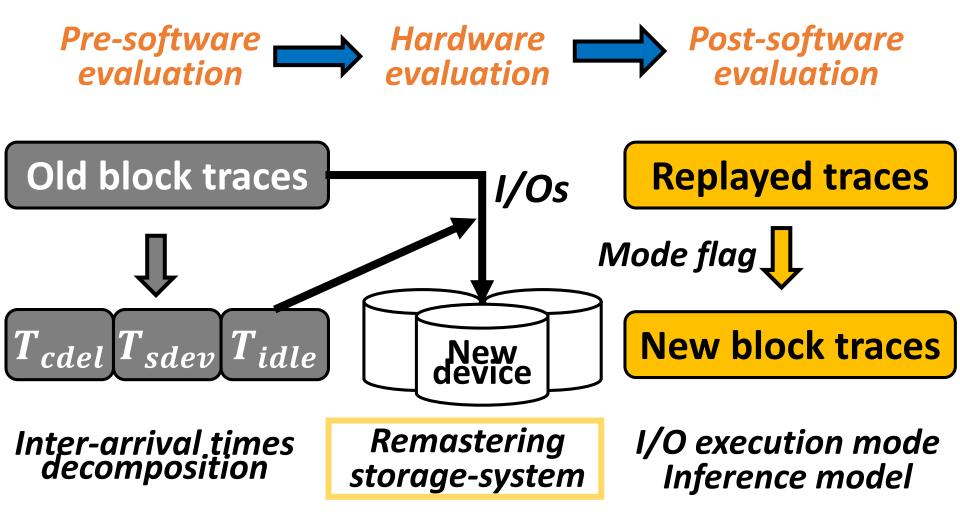




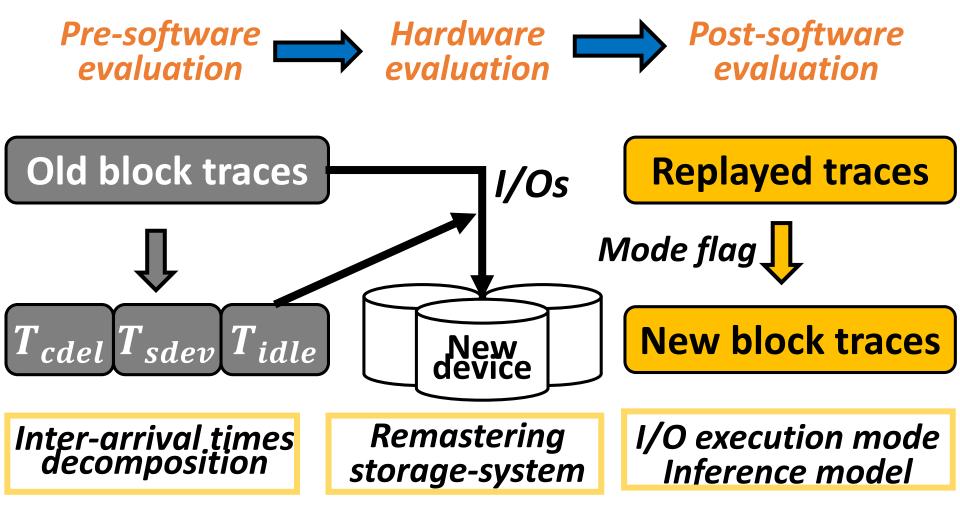
Inter-arrival times decomposition

Remastering storage-system I/O execution mode Inference model

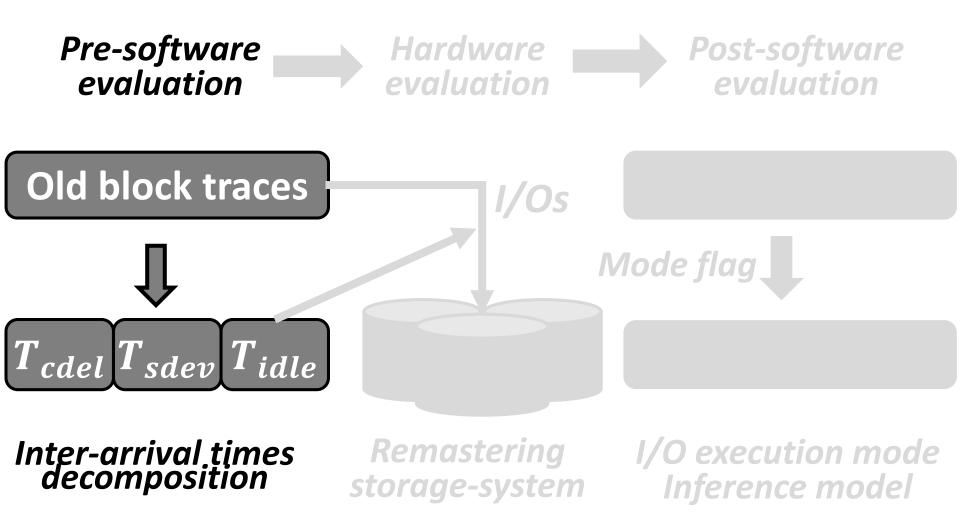














Request classification

Key idea: Classify all requests in trace with access pattern, operation type, request length

Old block traces

Req1: Read Len1 ..

Req2: Read Len2 ..

```
Req3: Write Len3 ..
```

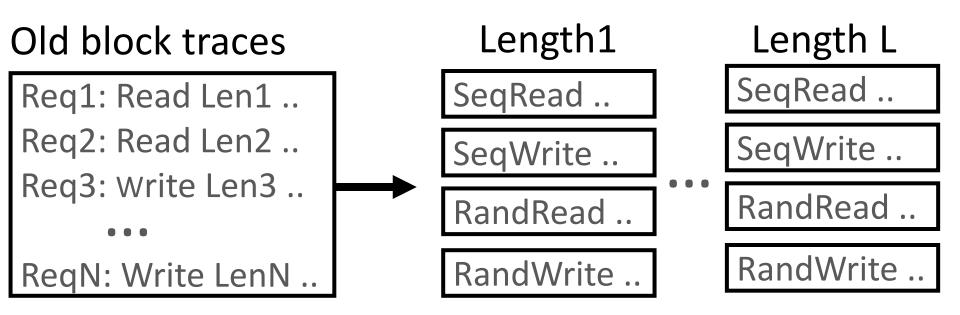
. . .

ReqN: Write LenN ..



Request classification

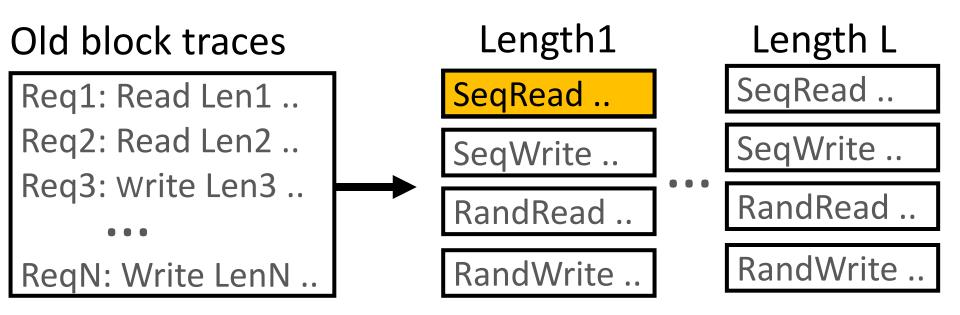
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Inter-arrival times distribution

Key idea: Easily decompose T_{intt} from same category requests' T_{intt} distribution



Key idea: Easily decompose T_{intt} from same category requests' T_{intt} distribution

Length1 sequential read

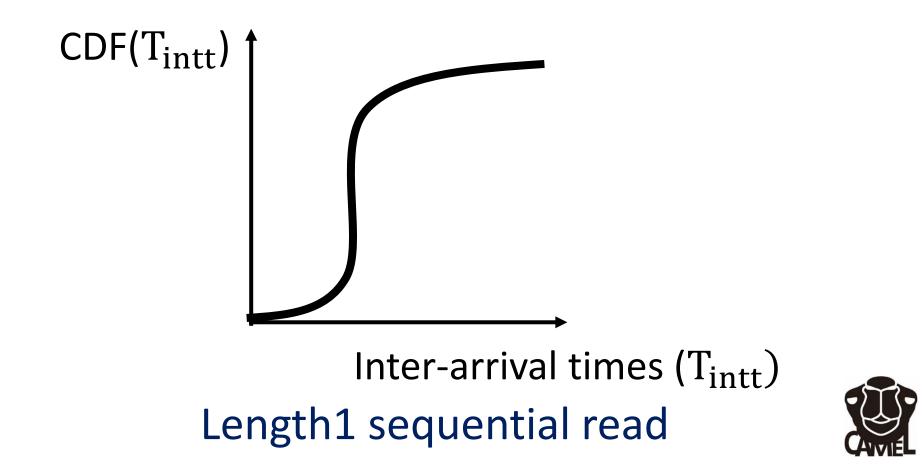


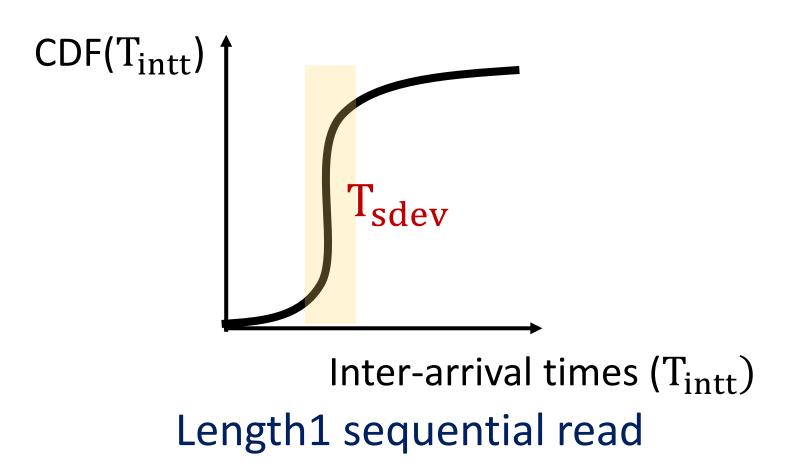
CDF(T_{intt})

Key idea: Easily decompose T_{intt} from same category requests' T_{intt} distribution

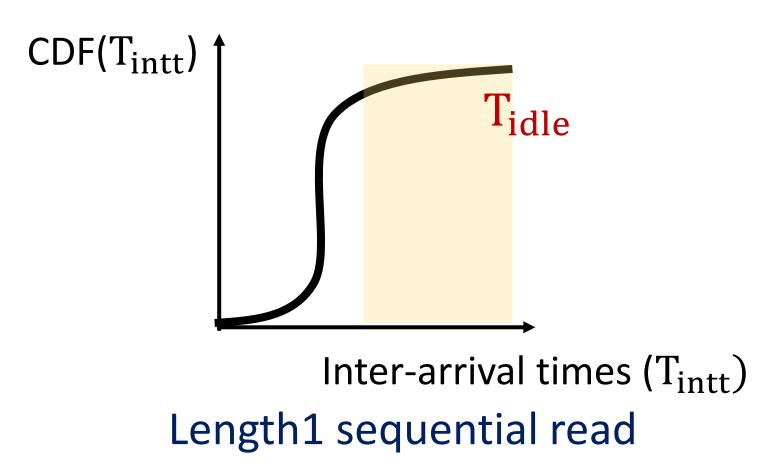
Inter-arrival times (T_{intt}) Length1 sequential read



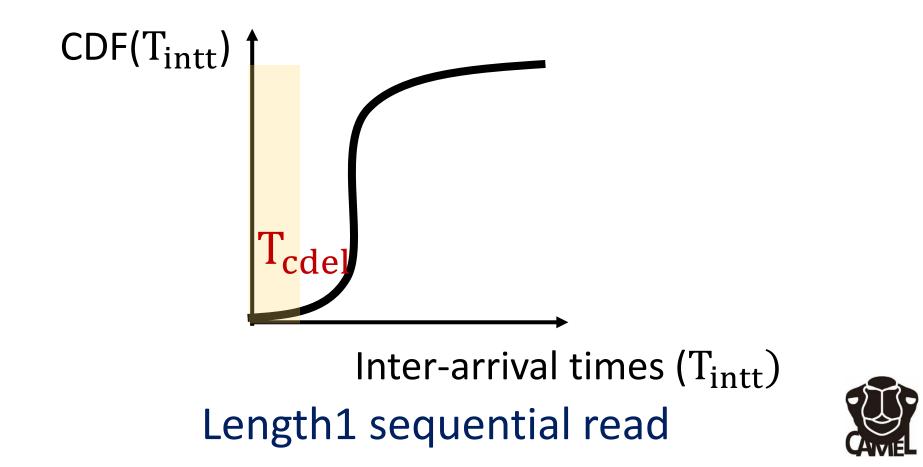




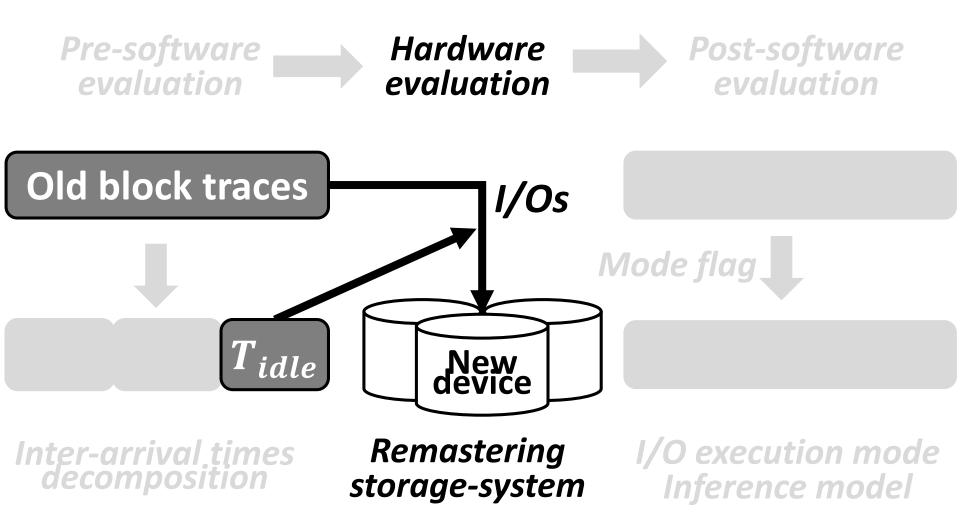








Overview of TraceTracker





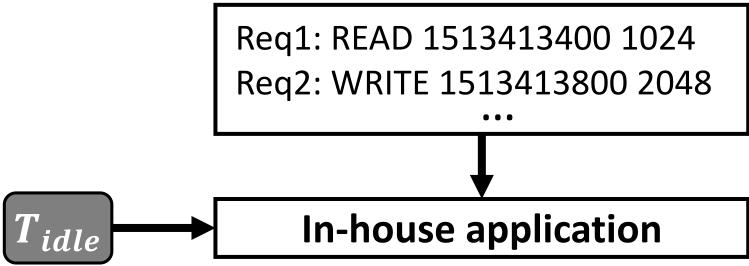
Trace replay

In-house application



Trace replay

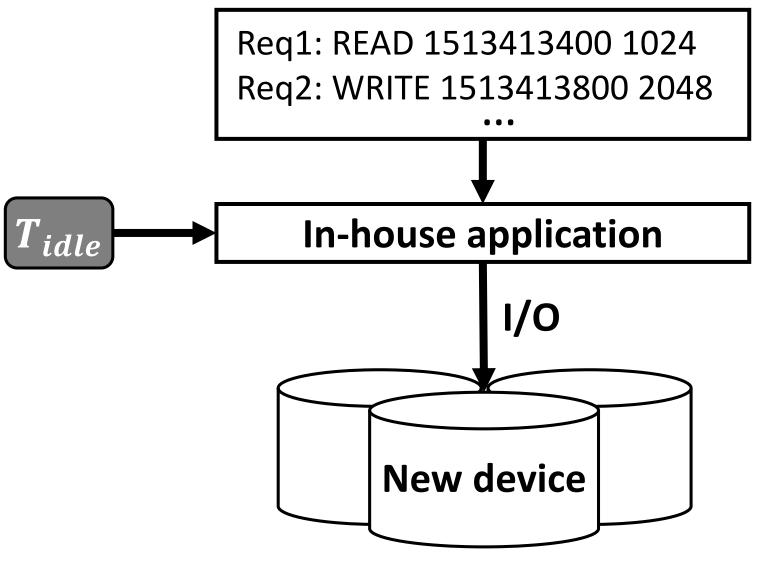
Old block trace





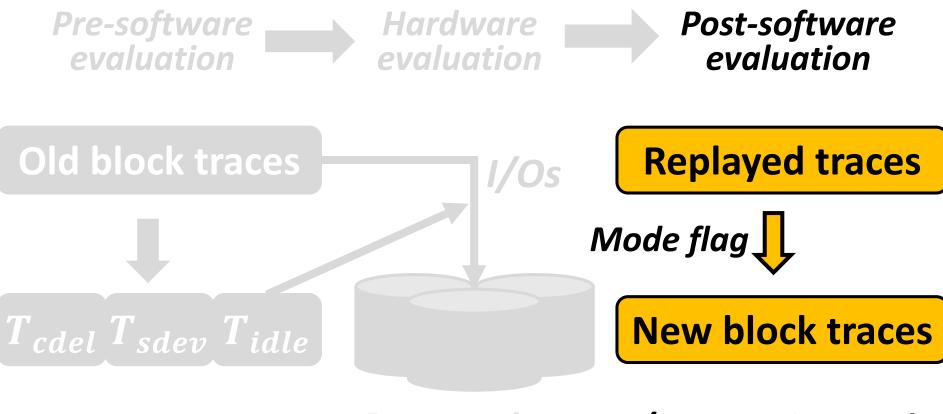
Trace replay

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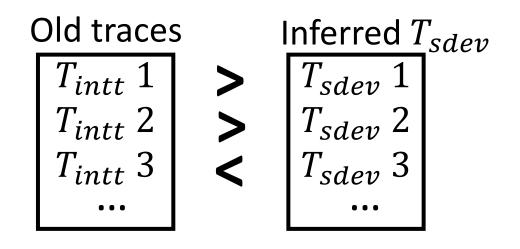
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Inter-arrival times decomposition Remastering storage-system I/O execution mode Inference model

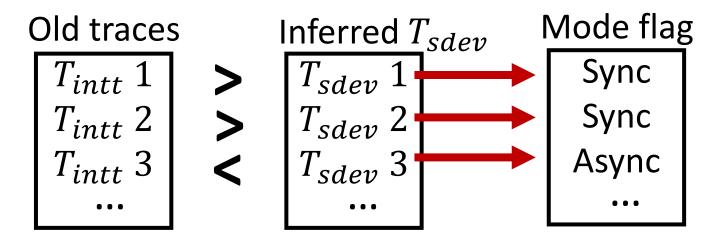


Key idea: Asynchronous I/O has short period which is less than storage I/O latency (T_{sdev})





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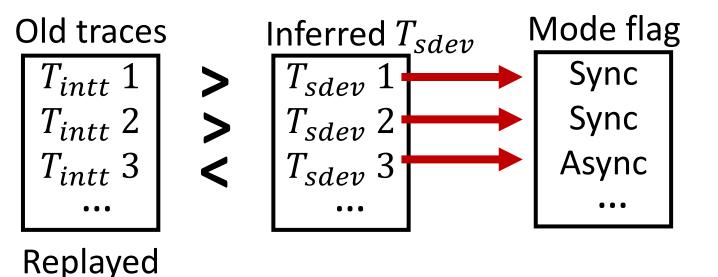


traces

 T_{intt} 1

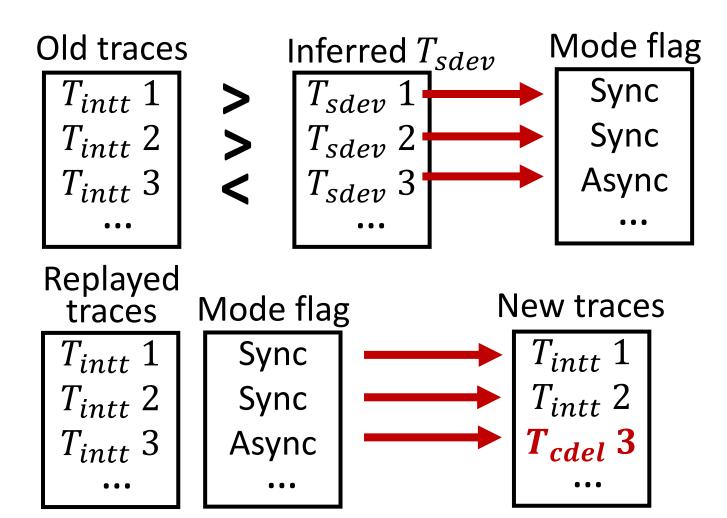
 T_{intt} 2

Key idea: Asynchronous I/O has short period which is less than storage I/O latency (T_{sdev})



AVVIE

Key idea: Asynchronous I/O has short period which is less than storage I/O latency (T_{sdev})





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Evaluation Methodology

Evaluation node

• All-flash array (by grouping four NVM Express SSDs)

Target block traces

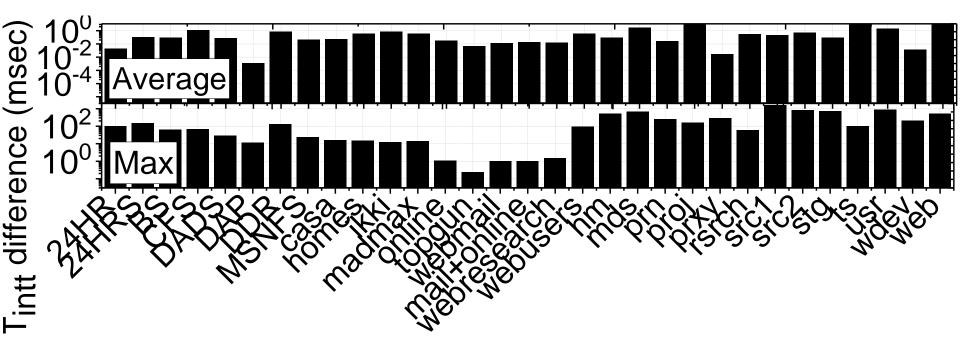
• 577 of large-scale block traces (around 2007~2009)

Reconstruction techniques

- Acceleration: Reconstruction by shortening T_{intt}
- Revision: Replaying block trace on all-flash array
- **Fixed-th**: An advanced revision method by inferring T_{idle} with a fixed threshold
- **Dynamic**: Reconstructions using our inference model, but with no post-processing.
- TraceTracker: Hardware/software co-evalution



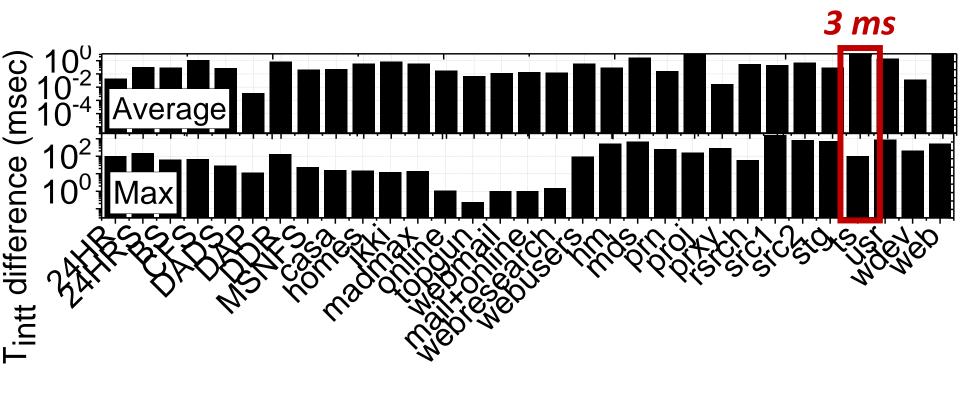
Evaluation #1: Different I/O timing



- New traces have, on average, 0.677 ms
 shorter inter-arrival time than old block traces
- Ex) 'ts' has 3 ms shorter T_{intt} , on average.



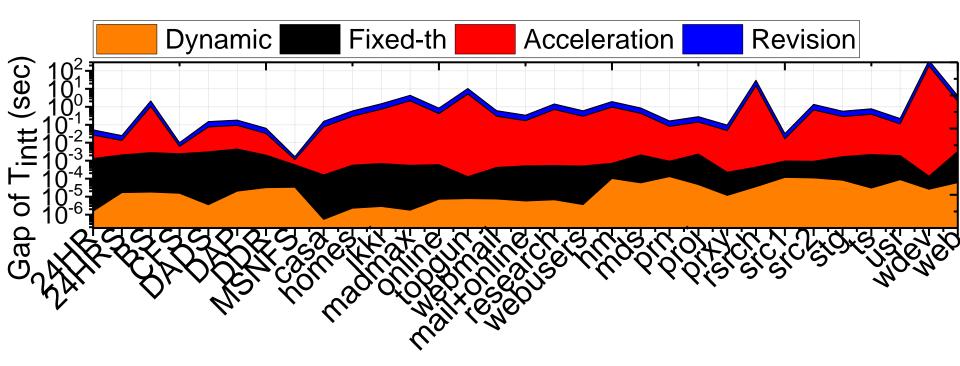
Evaluation #1: Different I/O timing



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Evaluation #2: Inaccuracy of reconstruction

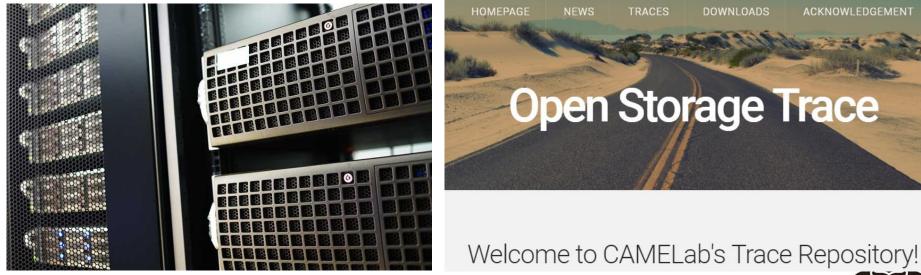


- Acceleration and Revision: 7.08, 7.15 sec due to no information for T_{idle}
- *Fixed-th* : 1.3 ms (inaccurate *T_{idle}*)
- Dynamic : 0.035 ms (inaccurate IO-mode)



Conclusion

- *TraceTracker* is trace reconstruction method which remasters the storage latency while maintaining runtime contexts of target traces.
- We're preparing brand new traces for open-license! Trace download is available at *trace.camelab.org*



KANDEMIR : All-Flash Array based HPC Testbed



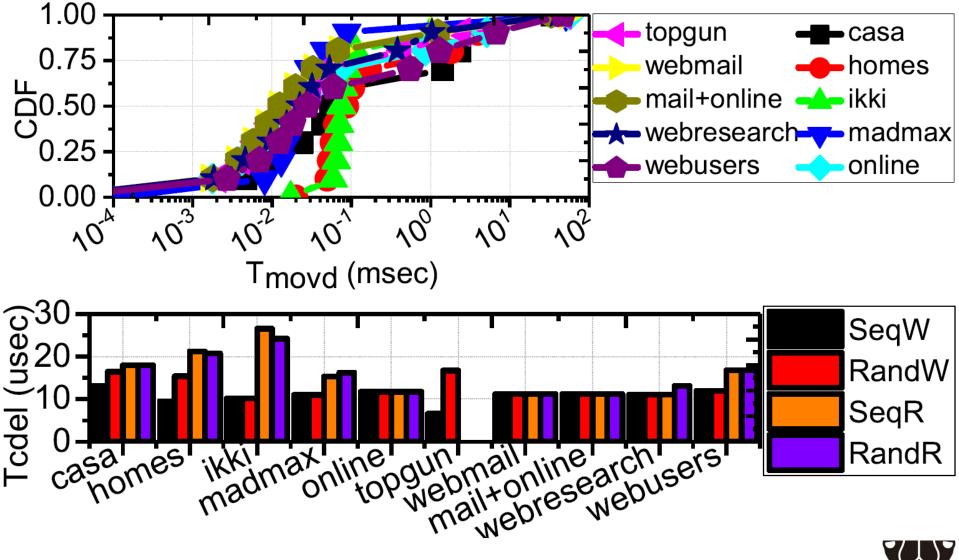
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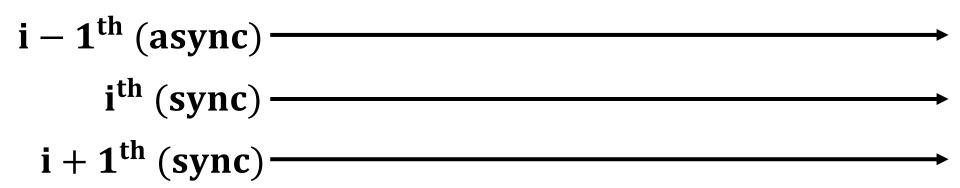
Miryeong Kwon, Jie Zhang, Gyuyoung Park, Wonil Choi, David Donofrio, John Shalf, Mahmut Kandemir, and Myoungsoo Jung



CDF distribution of *Tmovd*, *Tcdel*





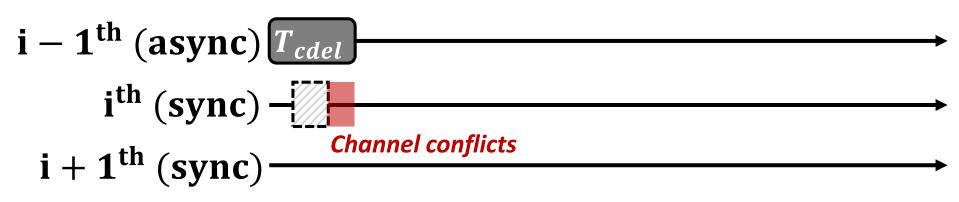




- $i 1^{th} (async) T_{cdel}$ $i^{th} (sync) - []_{-}$
 - $i + 1^{th} (sync) -$

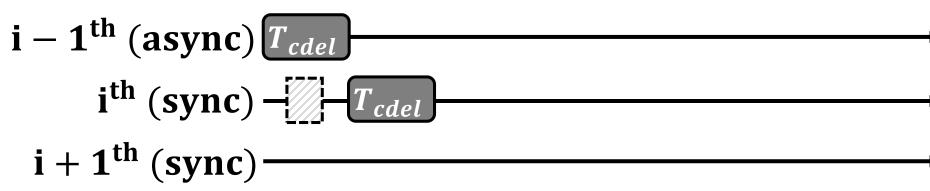
: Computation cycles by user/kernel





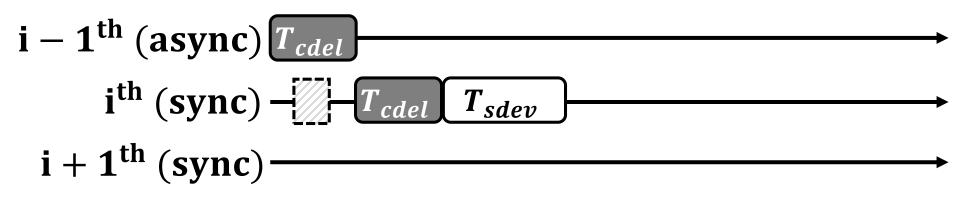
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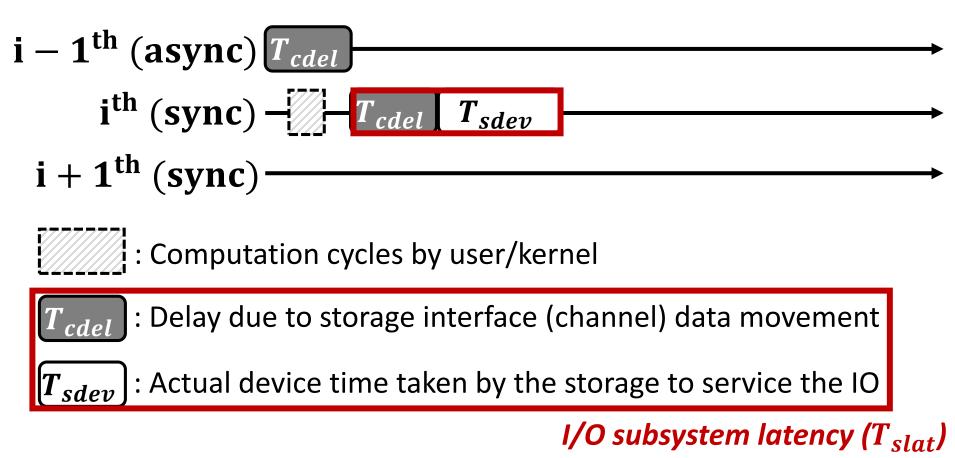




: Computation cycles by user/kernel

- tel : Delay due to storage interface (channel) data movement
- T_{sdev} : Actual device time taken by the storage to service the IO







 $i - 1^{th} (async) T_{cdel}$ $i^{th}(sync) - T_{cdel} T_{sdev}$ T_{idle} $i + 1^{th} (sync) -$

: Computation cycles by user/kernel



- $|T_{sdev}|$: Actual device time taken by the storage to service the IO
- : user/application does nothing T_{idle}



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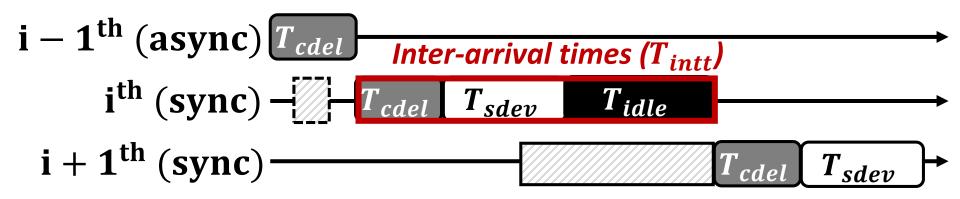
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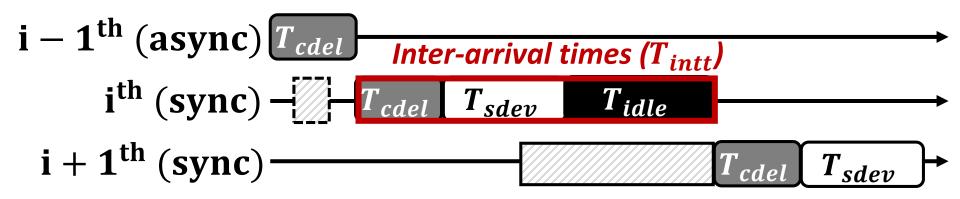


: Computation cycles by user/kernel

- **T**_{cdel}: Delay due to storage interface (channel) data movement
- T_{sdev} : Actual device time taken by the storage to service the IO
- Tidle : user/application does nothing

Inter-arrival times can be decomposed into T_{cdel}, T_{sdev}, and T_{idle}





: Computation cycles by user/kernel

- **T**_{cdel}: Delay due to storage interface (channel) data movement
- T_{sdev} : Actual device time taken by the storage to service the IO
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Why inter-arrival times decomposition is important to infer runtime context?



Overview of TraceTracker



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Pre-software evaluation



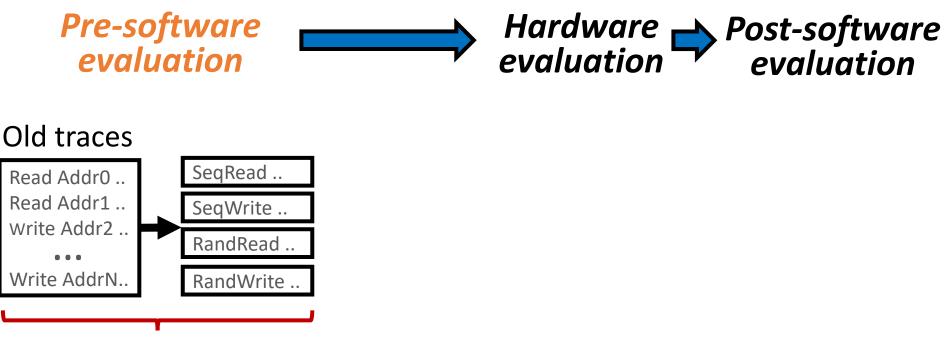
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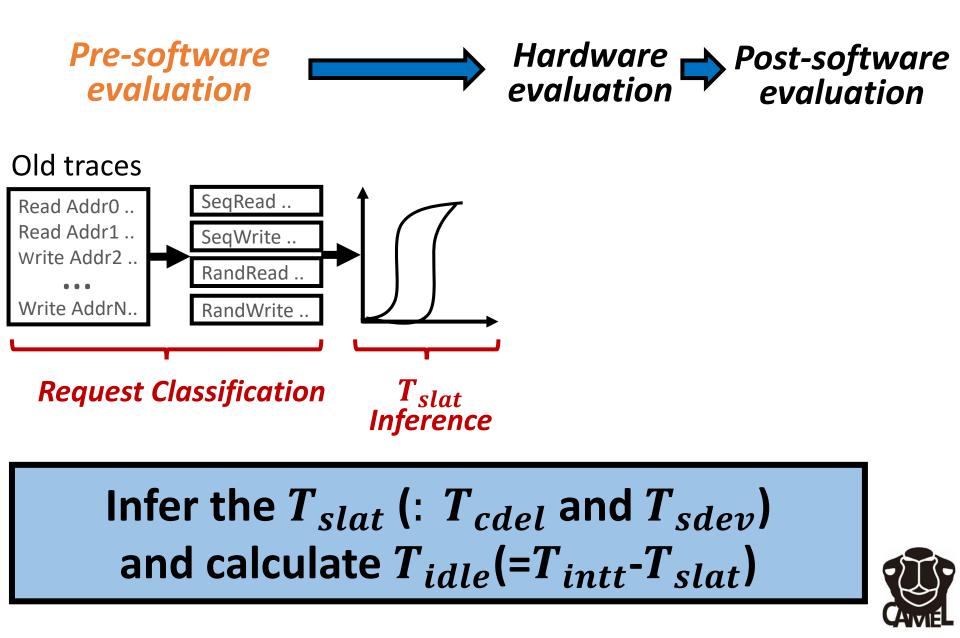


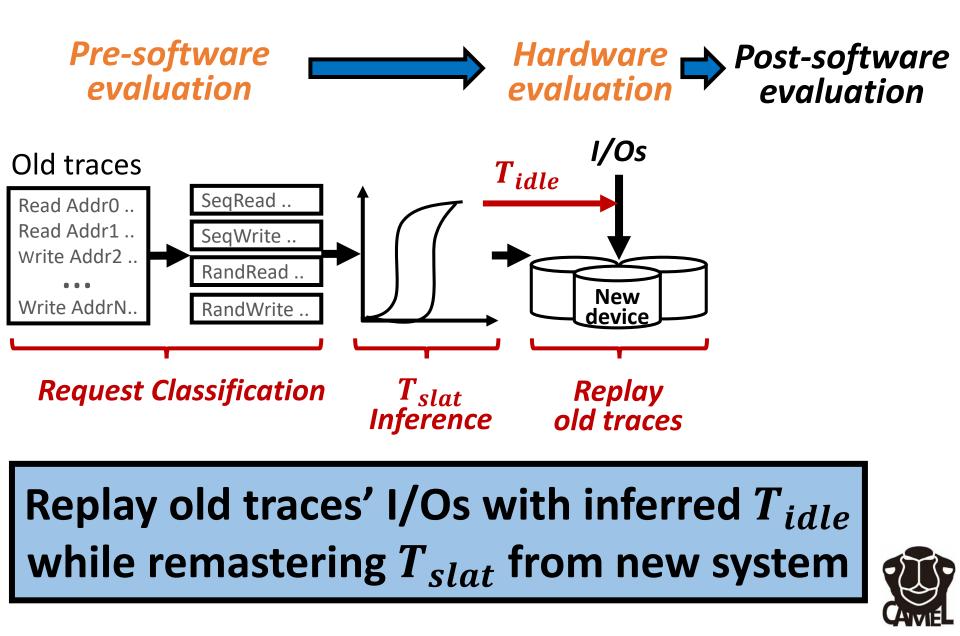


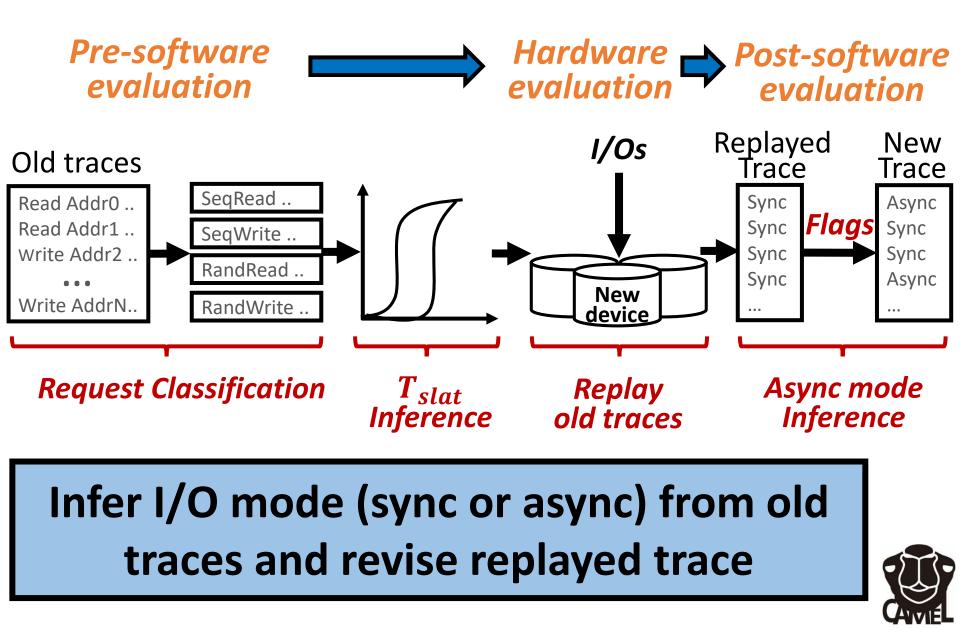
Request Classification

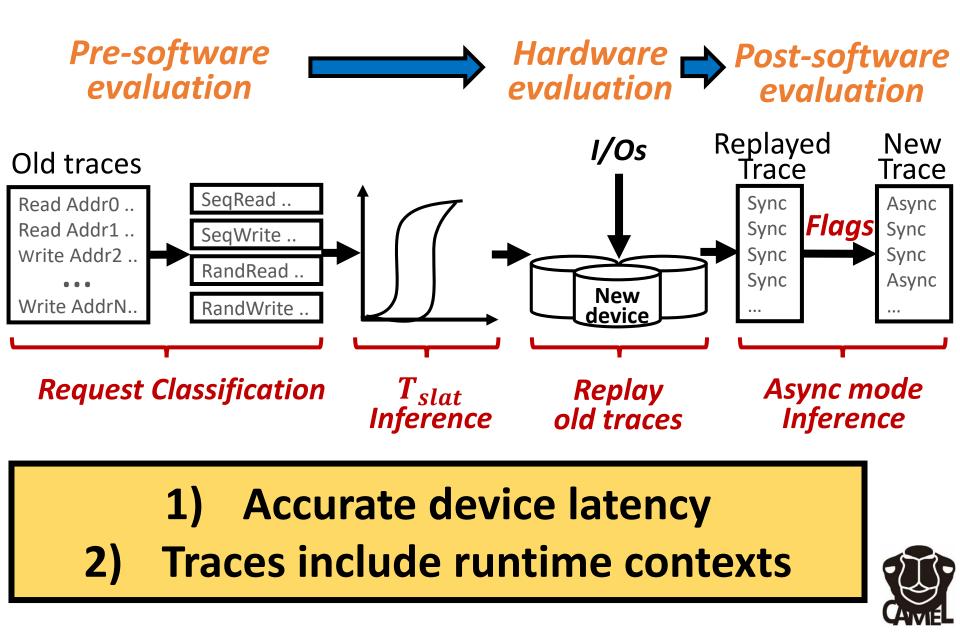
Each T_{cdel} and T_{sdev} have similar values for same operation type and request size

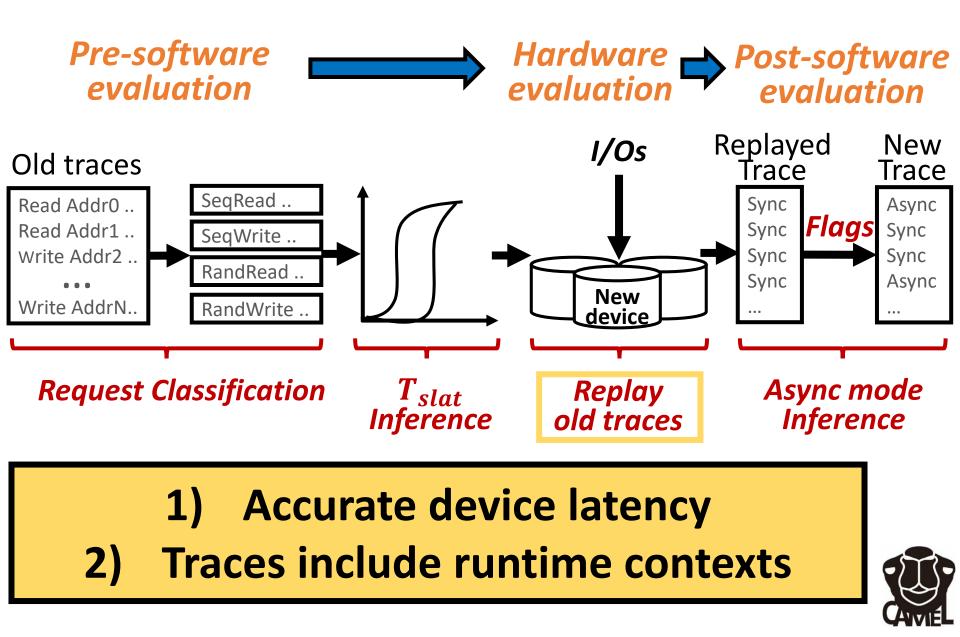


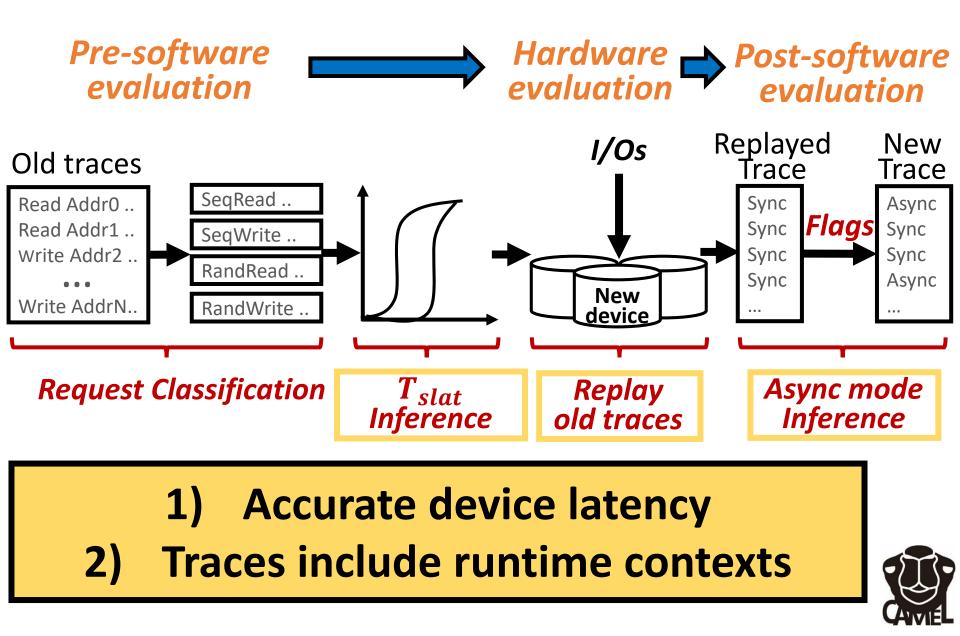












Key idea:

- 1) Each request group has similar T_{slat}
- 2) T_{sdev} can be expressed as linear model



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Sequential/Random Read

$$T_{slat} = T_{cdel_{read}} + Coeff_{read} \times REQSIZE + T_{movd}$$

Sequential/Random Write

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 $T_{slat}1 - T_{slat}2 = Coeff_{read} \times \frac{Diff(SIZE1 - SIZE2)}{Inferred from max slope of CDF}$ Known info.



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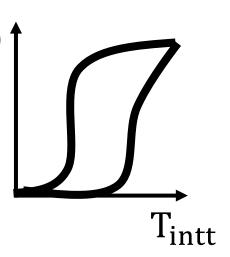
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 $CDF(T_{intt})$





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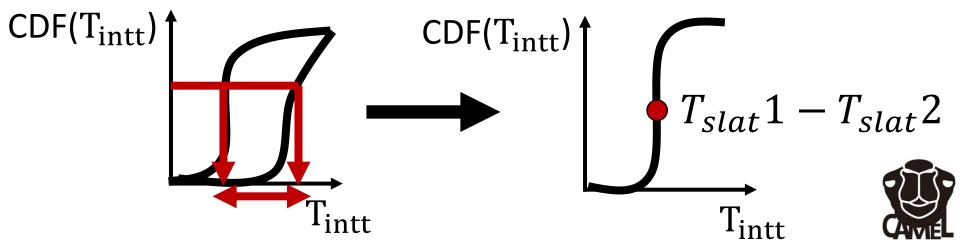
 $CDF(T_{intt}) \longrightarrow CDF(T_{intt}) \longrightarrow T_{intt}$

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Known info.

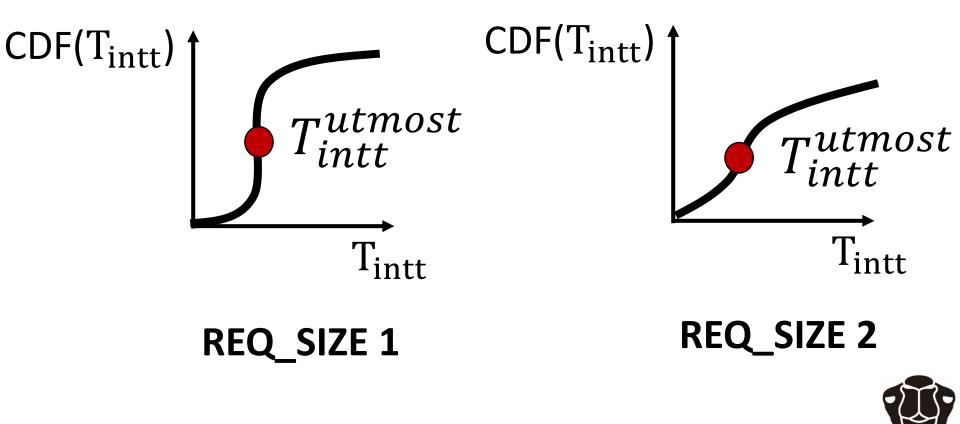
Then, which two CDF will be best?

Key idea: Select two steepest CDF



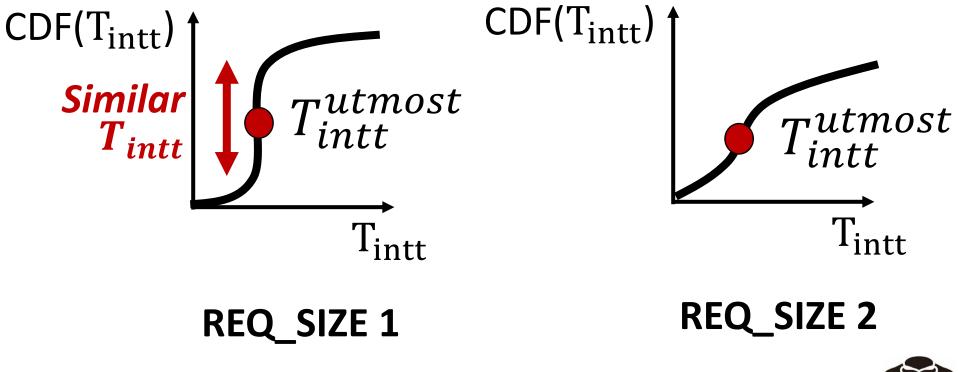
Key idea: Select two steepest CDF

 \because Steep CDF is more accurate to infer T_{slat}



Key idea: Select two steepest CDF

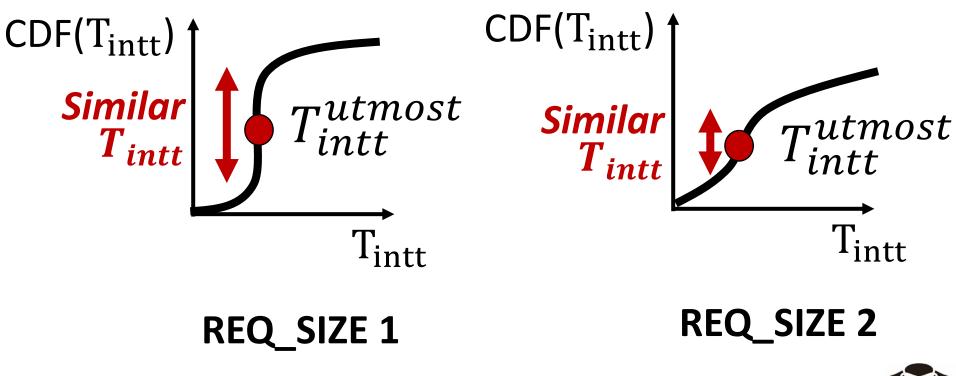
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Key idea: Mathematical analysis;

Maximum slope of CDF => Maximum of CDF'



Key idea: Mathematical analysis; Maximum slope of CDF => Maximum of CDF'

Challenge: Non-derivative discrete CDF data Solution: Use pchip-interpolation

Pchip interpolation



Key idea: Mathematical analysis; Maximum slope of CDF => Maximum of CDF'

Challenge: Non-derivative discrete CDF data Solution: Use pchip-interpolation

Pchip interpolation

Challenge: High cost of interpolation Solution: Use both PDF and CDF method



- **Key idea:** Mathematical analysis; Maximum slope of CDF => Maximum of CDF'
- Challenge: Non-derivative discrete CDF data Solution: Use pchip-interpolation
- Challenge: High cost of interpolation Solution: Use both PDF and CDF method

CDF(T_{intt}) *Tutmost CDF method*



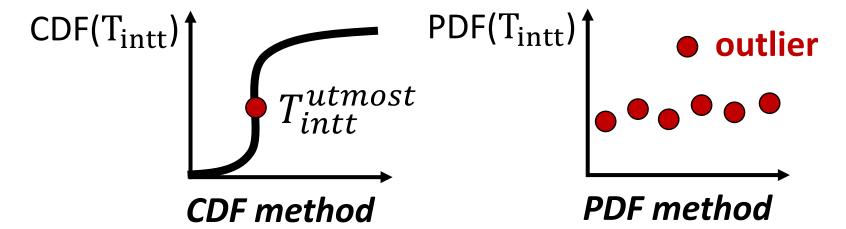
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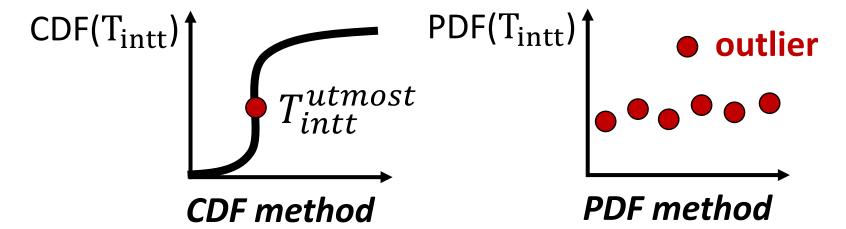


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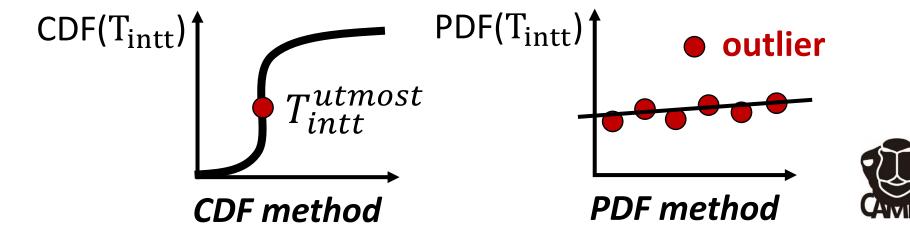
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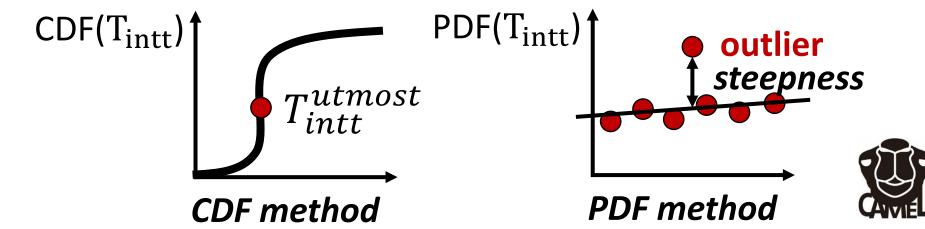


Pchip interpolation

Key idea: Mathematical analysis; Maximum slope of CDF => Maximum of CDF'

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Pchip interpolation

Solution:

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Pchip interpolation

Challenge: High cost of interpolation

	PDF method	CDF method
Accuracy	<	
Efficiency	>	



Solution:

Key idea: Mathematical analysis;

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	PDF method	CDF method
Accuracy	<	
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Steepness analysis



Solution:

Key idea: Mathematical analysis;

- Maximum slope of CDF => Maximum of CDF'
- Challenge: Non-derivative discrete CDF data Solution: Use pchip-interpolation

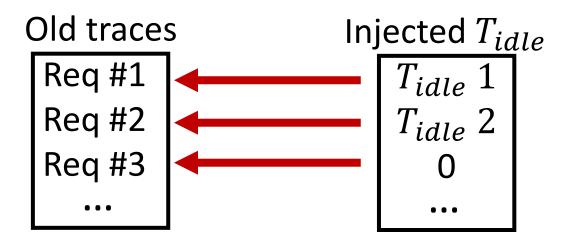
Pchip interpolation

Challenge: High cost of interpolation

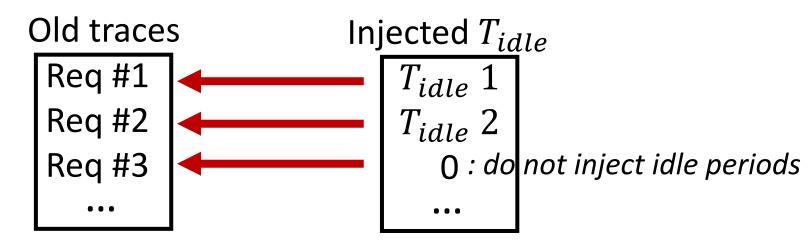
	PDF method	CDF method
Accuracy	<	
Efficiency	>	

Find max slope of CDF(diff

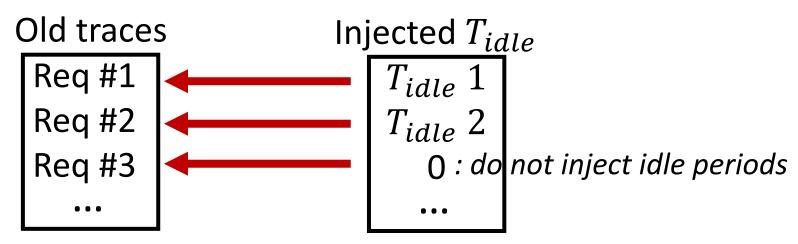




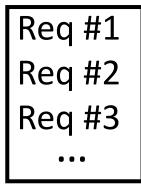




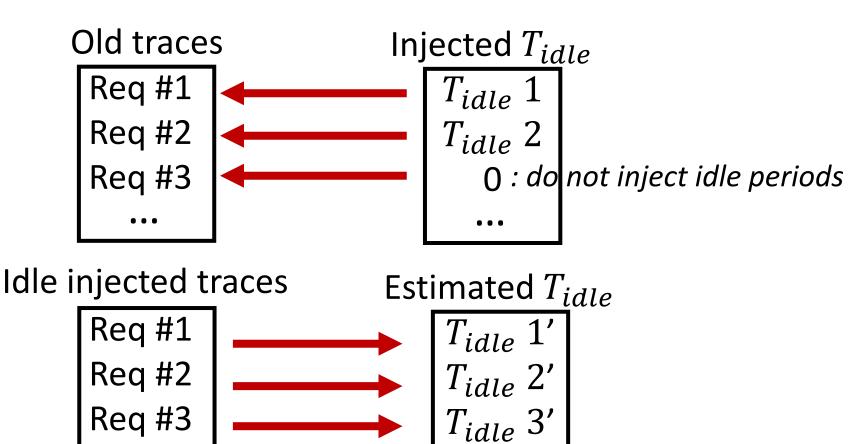




Idle injected traces

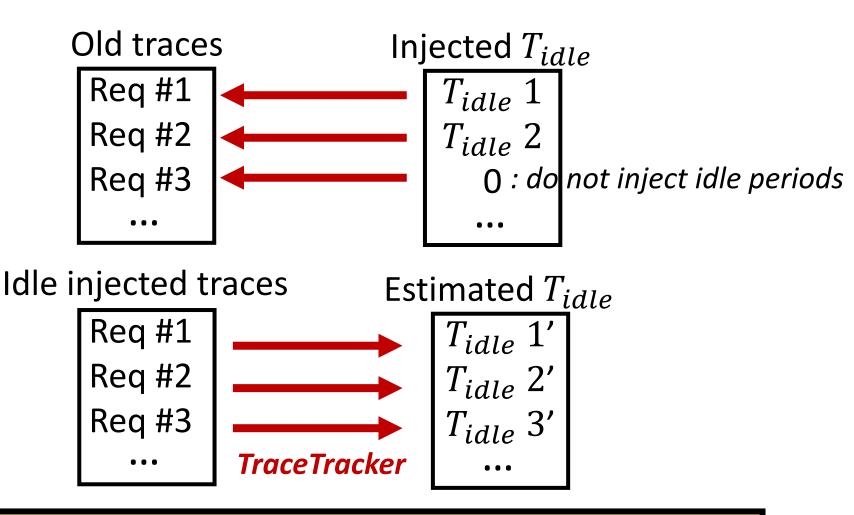






TraceTracker





Number of T_{idle} occurrence: 99%
 Total periods of T_{idle}: 96%

